A Paradox of Blown Leads: Rethinking Win Probability in Team Sports

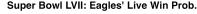
Jonathan Pipping and Abraham J. Wyner

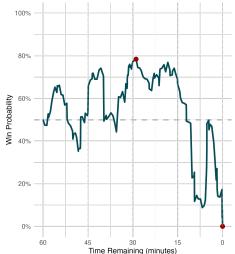
The Wharton School, University of Pennsylvania

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Case Study: Super Bowl LVII

- At the start of the 2nd half, the Chiefs trailed the Eagles 14–24 and faced a 3rd and 1 at their own 34
- At that point, the Eagles' projected win probability was 78.4%
- On the next play, Jerick McKinnon converted with a 14-yard run, and the Chiefs would go on to score, eventually winning 38–35.





Data via nflfastR

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- To investigate this, we formalize the question mathematically.

• The win probability of a team i at time t is a function of team strength S_i and game state (\mathcal{G}, t) . Symbolically,

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• We will investigate the distribution of M_{ℓ} as a function of the team strengths S_A and S_B .

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- ullet So the win prob. for each possible game state (\mathcal{G},t) is given by

$$\mathsf{WP}_A(t) = \mathbb{P}\left[\mathsf{Binom}(n_t, p_A) + \mathsf{score}_A(t) > \mathsf{Binom}(n_t, p_B) + \mathsf{score}_B(t)\right] + \frac{p_A}{p_A + p_B} \mathbb{P}(\mathsf{A} \text{ and } \mathsf{B} \text{ tie})$$

where $n_t = N - t$ is the number of possessions remaining.

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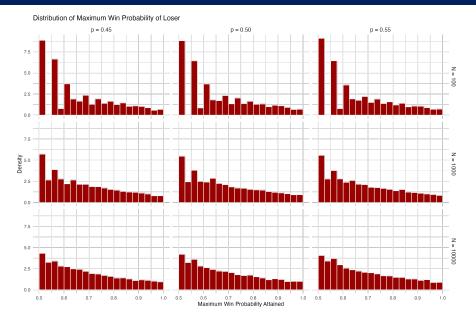
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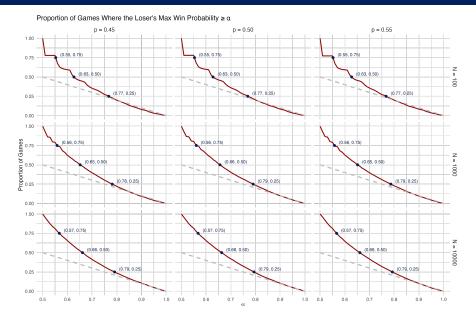
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 - Case 1: $p_A = p_B$, f is symmetric across teams
 - Case 2: $p_A \neq p_B$, f is asymmetric across teams

Symmetric Case: $p_A = p_B$



Symmetric Case: Threshold Plot



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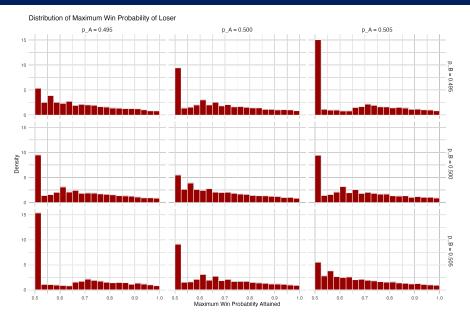
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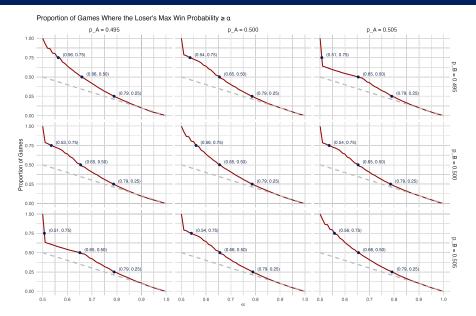
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- Holding $p_A = p_B$ constant, the distribution of M_ℓ becomes less discrete and approaches a continuous limit as N increases.
- In about half of all games, the losing team attains a win probability of at least 66% or more.

Asymmetric Case: $p_A \neq p_B$



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Asymmetric Case: Takeaways

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- Larger strength differentials decrease the proportion of games where the losing team attains a high win probability.

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- Then we can define Team A's win probability at each time step k as a Doob martingale:

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- Since τ_x is a stopping time, we invoke the optional stopping theorem to derive the distribution of M_A .

• **Theorem 1:** The distribution of M_A satisfies:

$$F_{M_A}(x) \ge 1 - \frac{p_0}{x}, \quad x \in [p_0, 1)$$

with equality when $\mathbb{P}(\tau_x = N) = 0$ (continuous limits).¹

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 Theorem 2: The conditional distribution of M_A given that team A loses satisfies:

$$F_{M_A|Y=0}(x) \ge 1 - \left(\frac{p_0}{1-p_0}\right) \cdot \left(\frac{1-x}{x}\right), \quad x \in [p_0, 1)$$

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- **Theorem 3:** Since team labels are arbitrary, let team A be the favorite $(p_0 \ge 0.5)$. Then the distribution of M_ℓ satisfies:

$$F_{M_{\ell}}(x) \geq egin{cases} 1 - rac{1 - p_0}{x} & ext{if } x \in [1 - p_0, p_0) \ 2 - rac{1}{x} & ext{if } x \in [p_0, 1) \end{cases}$$

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 - In this limit, the process becomes continuous and we can derive exact closed-form expressions for the distributions of M_A , M_B , and M_ℓ .
 - In addition, **Donsker's Invariance Principle** allows us to approximate the scoring process X_k with a Brownian motion!

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 - $\sigma^2 = p_A(1 p_A) + p_B(1 p_B)$ is the variance of each step in the scoring process.
- Then by Donsker's Invariance Principle, we have that

$$rac{X_{\lfloor Nt \rfloor}}{\sigma \sqrt{N}} \stackrel{d}{
ightarrow} B_t + \mu^* t, \quad \mu^* = rac{\sqrt{N} \, \mu}{\sigma}, \, \, t \in [0,1]$$

where B_t is standard Brownian motion.

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 - $\tau_x = \inf\{t \in [0,1] : p_t = x\}$ is the first time p_t exceeds x.
- Then we invoke the optional stopping theorem and use properties of Brownian motion to derive the distribution of M_A , M_B , and M_ℓ .

• **Theorem 4:** The distribution of M_A satisfies:

$$F_{M_A}(x) = 1 - \frac{p_0}{x} = 1 - \frac{\Phi(\mu^*)}{x}, \quad x \in [\Phi(\mu^*), 1)$$

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 Theorem 5: The conditional distribution of M_A given that team A loses satisfies:

$$F_{M_A|Y=0}(x) = 1 - \left(\frac{p_0}{1 - p_0}\right) \cdot \left(\frac{1 - x}{x}\right), \quad x \in [p_0, 1)$$
$$= 1 - \left(\frac{\Phi(\mu^*)}{1 - \Phi(\mu^*)}\right) \cdot \left(\frac{1 - x}{x}\right), \quad x \in [\Phi(\mu^*), 1)$$

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- We saw in the discrete case that the distribution of M_{ℓ} is a piecewise mixture. We use a similar logic to derive the distribution of M_{ℓ} .
- Theorem 6: Since team labels are arbitrary, let team A be the favorite $(\mu^* \geq 0)$. Then the distribution of M_ℓ satisfies

$$F_{M_{\ell}}(x) = egin{cases} 1 - rac{1 - p_0}{x} & ext{if } x \in [1 - p_0, p_0) \ 2 - rac{1}{x} & ext{if } x \in [p_0, 1) \end{cases}$$

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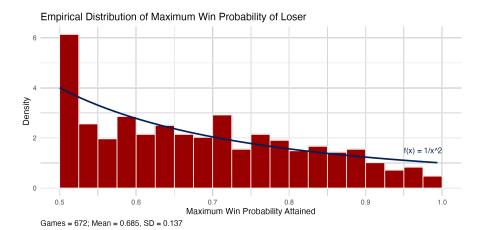
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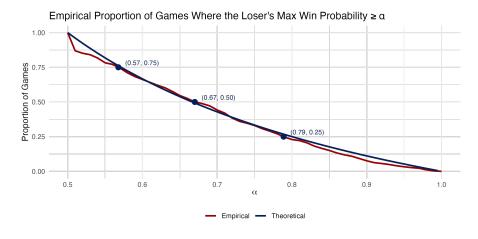
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 - Scoring often isn't binary!
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 - Weather, injuries, game strategy, and other external factors matter!
- Let's take a look at the distribution of M_{ℓ} for real-life games. . .

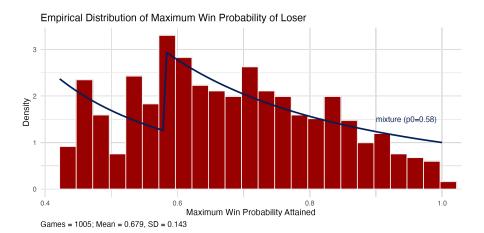
Evenly-Matched NFL Games (2002–2024, |Spread| < 2)



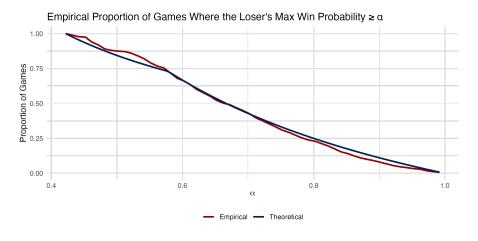
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Unevenly-Matched NFL Games: 2002-2024, |Spread| = 3



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- Blown leads happen all the time, especially in even matchups!
- Conditioning on an eventual loss fundamentally changes the distribution of the maximum win probability attained.

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Conclusions

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Conclusions

- In Super Bowl LVII, the Eagles reached a maximum win probability of 78.4% before ultimately losing.
- From this position (assuming a well-calibrated model) **it's true** that the Eagles only had a 21.6% chance of losing the game.
- However, the event that the eventual loser of this game reached 78.4% is **provably closer** to 30%.

Acknowledgements

- Special thanks to Professor Jiaoyang Huang, Dr. Paul Sabin, and Dr. Ryan Brill for their helpful feedback.
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• **Theorem 1:** When $p_A = p_B$, the cumulative distribution of M satisfies

$$F_M(x) \ge 1 - \frac{p_0}{x}, \quad x \in [p_0, 1)$$

with equality exactly when $\mathbb{P}(\tau_{\mathsf{x}} = \mathsf{N}) = 0$.

Proof of Theorem 1 (continued)

- If $x \le p_0$, then $\tau_x = 0$ a.s., so $M \ge x$ a.s. and $F_M(x) = 0$.
- For $x > p_0$, by the optional stopping theorem on the bounded martingale (p_k) at $\tau_x \wedge N$:

$$\mathbb{E}[p_{\tau_{\times} \wedge N}] = p_0$$

Decomposing this:

$$p_0 = \mathbb{E}[p_{\tau_x} \mathbf{1}_{\{\tau_x < N\}}] + \mathbb{E}[p_N \mathbf{1}_{\{\tau_x = N\}}]$$

• Since $p_{\tau_x} = x$ on $\{\tau_x < N\}$ and $p_N = 1$ on $\{\tau_x = N\}$:

$$p_0 = x \mathbb{P}(\tau_x < N) + \mathbb{P}(\tau_x = N)$$

Proof of Theorem 1 (continued)

• Since $\mathbb{P}(M \ge x) = \mathbb{P}(\tau_x < N) + \mathbb{P}(\tau_x = N)$:

$$p_0 = x \mathbb{P}(M \ge x) + (1 - x) \mathbb{P}(\tau_x = N)$$

• For x < 1, $\mathbb{P}(\tau_x = N) \ge 0$ and we have:

$$\mathbb{P}(M \geq x) \leq \frac{p_0}{x}, \quad x \in [p_0, 1)$$

with equality exactly when $\mathbb{P}(\tau_x = N) = 0$.

• When x=1, $\tau_x=N$ a.s., so:

$$\mathbb{P}(M \geq 1) = \mathbb{P}(M = 1) = p_0$$

• **Theorem 2:** When $p_A = p_B$, the cumulative distribution of M conditional on Y = 0 satisfies

$$F_{M|Y=0}(x) \ge 1 - \left(\frac{p_0}{1-p_0}\right) \cdot \left(\frac{1-x}{x}\right), \quad x \in [p_0, 1)$$

with equality exactly when $\mathbb{P}(\tau_x = N) = 0$.

Proof of Theorem 2 (continued)

• Decompose the event $\{M \ge x\}$:

$$\mathbb{P}(M \ge x) = \mathbb{P}(M \ge x, Y = 0) + \mathbb{P}(M \ge x, Y = 1)$$

• On $\{Y = 1\}$, $p_N = 1 \ge x$ a.s. So $\{M \ge x\} \cap \{Y = 1\} = \{Y = 1\}$:

$$\mathbb{P}(M \ge x) = \mathbb{P}(Y = 1) + \mathbb{P}(M \ge x, Y = 0)
= \mathbb{P}(Y = 1) + \mathbb{P}(Y = 0)\mathbb{P}(M \ge x \mid Y = 0)
= \rho_0 + (1 - \rho_0)\mathbb{P}(M \ge x \mid Y = 0)$$

- From Theorem 1: $\frac{p_0}{x} \ge p_0 + (1-p_0)\mathbb{P}(M \ge x \mid Y=0)$
- Solving for $\mathbb{P}(M \ge x \mid Y = 0)$:

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with equality exactly when $\mathbb{P}(\tau_x = N) = 0$.

• **Theorem 3:** When $p_A = p_B$, the distribution of M_ℓ satisfies

$$F_{M_{\ell}}(x) \ge 2 - \frac{1}{x}, \quad x \in [0.5, 1)$$

with equality exactly when $\mathbb{P}(\tau_x = N) = 0$.

Proof of Theorem 3 (continued)

- Since $p_A = p_B$, we have $p_0 = 0.5$ and the mixture weights are equal.
- From Theorem 2, both conditional distributions are identical:

$$F_{M_A|Y=0}(x) = F_{M_B|Y=1}(x) \ge 1 - \left(\frac{0.5}{0.5}\right) \cdot \left(\frac{1-x}{x}\right) = 1 - \frac{1-x}{x}$$

• Since both teams have identical distributions, the mixture is simply:

$$F_{M_{\ell}}(x) \ge 0.5 \cdot F_{M_A|Y=0}(x) + 0.5 \cdot F_{M_B|Y=1}(x) = 1 - \frac{1-x}{x} = 2 - \frac{1}{x}$$

• **Theorem 4:** When $p_A \neq p_B$, the cumulative distribution of M satisfies

$$F_M(x) = 1 - \frac{p_0}{x} = 1 - \frac{\Phi(\mu^*)}{x}, \quad x \in [p_0, 1)$$

Proof of Theorem 4 (continued)

• By the optional stopping theorem on the bounded martingale (p_t) at $\tau_x \wedge 1$:

$$\mathbb{E}[p_{\tau_{\mathsf{x}} \wedge 1}] = p_0$$

Decomposing this:

$$p_0 = \mathbb{E}[p_{\tau_x} \mathbf{1}_{\{\tau_x < 1\}}] + \mathbb{E}[p_1 \mathbf{1}_{\{\tau_x = 1\}}]$$

• Since $p_{\tau_x} = x$ on $\{\tau_x < 1\}$ and $p_1 = 1$ on $\{\tau_x = 1\}$:

$$p_0 = x \mathbb{P}(\tau_x < 1) + \mathbb{P}(\tau_x = 1)$$

• Since $\mathbb{P}(M \ge x) = \mathbb{P}(\tau_x < 1) + \mathbb{P}(\tau_x = 1)$:

$$p_0 = x \mathbb{P}(M \ge x) + \mathbb{P}(\tau_x = 1)$$

Proof of Theorem 4 (continued)

• For x < 1, $\mathbb{P}(\tau_x = 1) = 0$ (continuous paths), so:

$$\mathbb{P}(M \ge x) = \frac{p_0}{x} = \frac{\Phi(\mu^*)}{x}, \quad x \in [p_0, 1)$$

• When x = 1, $\tau_x = 1$ a.s., so:

$$\mathbb{P}(M \geq 1) = \mathbb{P}(M = 1) = \rho_0 = \Phi(\mu^*)$$

• **Theorem 5:** When $p_A \neq p_B$, the conditional distribution of M given Y = 0 satisfies

$$F_{M|Y=0}(x) = 1 - \left(\frac{p_0}{1 - p_0}\right) \cdot \left(\frac{1 - x}{x}\right), \quad x \in [p_0, 1)$$

where $p_0 = \Phi(\mu^*)$ and M = 0 for $x \in [0, p_0)$.

Proof of Theorem 5 (continued)

• Decompose the event $\{M \ge x\}$:

$$\mathbb{P}(M \ge x) = \mathbb{P}(M \ge x, Y = 0) + \mathbb{P}(M \ge x, Y = 1)$$

• On $\{Y = 1\}$, $p_1 = 1 \ge x$ a.s. So $\{M \ge x\} \cap \{Y = 1\} = \{Y = 1\}$:

$$\mathbb{P}(M \ge x) = \mathbb{P}(Y = 1) + \mathbb{P}(M \ge x, Y = 0)
= \mathbb{P}(Y = 1) + \mathbb{P}(Y = 0)\mathbb{P}(M \ge x \mid Y = 0)
= p_0 + (1 - p_0)\mathbb{P}(M \ge x \mid Y = 0)$$

- From Theorem 4: $\frac{p_0}{x} = p_0 + (1 p_0) \mathbb{P}(M \ge x \mid Y = 0)$
- Solving for $\mathbb{P}(M \ge x \mid Y = 0)$:

$$\mathbb{P}(M \ge x \mid Y = 0) = \left(\frac{p_0}{1 - p_0}\right) \cdot \left(\frac{1 - x}{x}\right)$$

• **Theorem 6:** When $p_A \neq p_B$, let team A be the favorite $(\mu^* \geq 0)$. Then the distribution of M_ℓ satisfies

$$F_{M_{\ell}}(x) = \begin{cases} 1 - \frac{1 - p_0}{x} & \text{if } x \in [1 - p_0, p_0) \\ 2 - \frac{1}{x} & \text{if } x \in [p_0, 1) \end{cases}$$

where $p_0 = \Phi(\mu^*)$ and $M_\ell = 0$ for $x \in [0, 1 - p_0)$.

Proof of Theorem 6 (continued)

- The distribution of M_{ℓ} is a mixture: $F_{M_{\ell}}(x) = (1 - p_0)F_{M_{\Delta}|Y=0}(x) + p_0F_{M_{B}|Y=1}(x).$
- From Theorem 5, we have the conditional distributions:

$$F_{M_A|Y=0}(x) = \begin{cases} 0 & x \in [0, p_0) \\ 1 - \left(\frac{p_0}{1-p_0}\right) \cdot \left(\frac{1-x}{x}\right) & x \in [p_0, 1) \end{cases}$$

$$F_{M_B|Y=1}(x) = \begin{cases} 0 & x \in [0, 1-p_0) \\ 1 - \left(\frac{1-p_0}{p_0}\right) \cdot \left(\frac{x}{1-x}\right) & x \in [1-p_0, 1) \end{cases}$$

Proof of Theorem 6 (continued)

• For $x \in [1 - p_0, p_0)$: $F_{M_A|Y=0}(x) = 0$ (since $x < p_0$), so

$$egin{split} F_{M_\ell}(x) &= (1-p_0)\cdot 0 + p_0\cdot \left(1-rac{1-p_0}{p_0}\cdot rac{x}{1-x}
ight) \ &= 1-rac{1-p_0}{x} \end{split}$$

• For $x \in [p_0, 1)$: Both terms contribute, giving

$$F_{M_{\ell}}(x) = (1 - p_0) \left(1 - \frac{p_0}{1 - p_0} \cdot \frac{1 - x}{x} \right) + p_0 \left(1 - \frac{1 - p_0}{p_0} \cdot \frac{x}{1 - x} \right)$$
$$= 2 - \frac{1}{x}$$